Talagrand's ellipsoid construction

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1

Gaussian processes indexed by ellipsoids

S:gaussian

The set $\ell^2 := \{s \in \mathbb{R}^{\mathbb{N}} : \sum_{n \in \mathbb{N}} s_n^2\}$ is a separable HILBERT space under the inner product $\langle s, t \rangle = \sum_{n \in \mathbb{N}} s_n t_n$, with corresponding norm $||s||_2 = (\sum_{n \in \mathbb{N}} s_n^2)^{1/2}$. If $\{g_n : n \in \mathbb{N}\}$ is a sequence of independent N(0, 1) random variables then

$$X_t := \sum_n t_n g_n \qquad \text{for } t \in \ell^2$$

converges in both the almost sure and L^2 senses. The collection $\{X_t : t \in \ell^2\}$ is a centered gaussian process with

$$\operatorname{cov}(X_s, X_t) = \sum_{\alpha} s_{\alpha} t_{\alpha} = \langle s, t \rangle \quad \text{for } s, t \in \ell^2.$$

Without loss of generality we may assume that the process is DOOB-SEPARABLE.

Talagrand (2021, pp. 73–80) considered the case where T is an ellipsoid, that is,

$$T = \mathcal{E}[a] := \{ t \in \ell^2 : \sum_n (t_n/a_n)^2 \le 1 \}$$
 for some fixed a in ℓ^2 .

The set T is a compact, convex subset of ℓ^2 . Talagrand noted (Prop. 2.13.1) that $\mathbb{P}\sup_{t\in T} X_t$ is of order $||a||_2$. His argument can be slightly simplified.

eea <1>

Lemma. For $T := \mathcal{E}[a]$,

$$c_0 ||a||_2 \le \mathbb{P} \sup_t X_t \le ||a||_2$$
 where $c_0 := \sqrt{2/\pi} = \mathbb{P}|g_1|$.

Proof. Rescale X_t as $\langle u, W \rangle$ where $W_i := a_i g_i$ and $u_i := t_i / a_i$, so that

$$\sup_{t \in T} X_t = \sup\{\langle u, W \rangle : \|u\|_2 \le 1\}$$

By CAUCHY-SCHWARZ, $\langle u, W \rangle \leq ||u||_2 ||W||_2 = ||W||_2$ with equality when $u_i = W_i / ||W||_2$. Thus $\sup_{t \in T} X_t = ||W||_2$ and

$$\mathbb{P}\sup_{t\in T} X_t = \mathbb{P} \|W\|_2 \le \sqrt{\mathbb{P} \|W\|_2^2} = \|a\|_2.$$

For the lower bound, let S denote the set of sequences $u = (u_1, u_2, ...)$ with $u_i = \pm a_i / ||a||_2$. Each such u has $||u||_2^2 = \sum_i a_i^2 / ||a||_2^2 = 1$ and

$$\sup_{u \in S} \langle u, W \rangle = \sup_{u \in A} \sum_{i} \pm a_i^2 g_i / \|a\|_2 = \sum_{i} a_i^2 |g_i| / \|a\|_2,$$

 $\square \quad \text{with expected value } \sum_{i} a_i^2 c_0 / \|a\|_2 = c_0 \|a\|_2.$

Talagrand (2021, pp. 73–80) showed how to construct subsets $T_n \subset \ell^2$ with

$$U_0 = \{0\} \text{ and } |U_n| \le N_n := 2^{2^n} \text{ and}$$
$$\sup_{t \in T} \sum_{n \ge 0} 2^{n/2} d(t, U_n) \le C_1 ||a||_2,$$

for some constant C_1 .

Remark. It suffices to construct the U_n 's so that $\log_2 |U_n| \leq 2^{n+c}|$ for some universal constant c. Then work with the approximating sets $\widetilde{U}_n = \{0\}$ for $n \leq m$ and $\widetilde{U}_n = U_{n+m}$ for some integer m greater than c. The subset \widetilde{U}_n could then be replaced by $\{\pi(s) : s \in \widetilde{U}_n\}$, where π denotes the map that takes each s in ℓ^2 to its closest point in T. For each s in \widetilde{U}_n and t in T we have $||s - t||_2 \geq ||s - \pi(s)||_2$, which implies

 $\|t - \pi(s)\|_{2} \le \|t - s\|_{2} + \|s - \pi(s)\|_{2} \le 2 \|t - s\|_{2}.$

S:blocks

 $\mathbf{2}$

Reduction to blocks

Without loss of generality we can assume that $a_i \searrow 0$. For the purpose of controlling the size of the approximating set U_n we will need to partition \mathbb{N} into disjoint blocks E_1, E_2, \ldots . My choice of blocks will differ slightly from the blocks used by Talagrand. For each positive integer n define

$$E_n := \{i \in \mathbb{N} : \alpha_n \le i < \alpha_{n+1}\} \quad \text{where } \alpha_n := 2^{n-1},$$
$$J_n := \{i \in \mathbb{N} : 1 \le i < \alpha_{n+1}\} = \{i \in \mathbb{N} : i < 2^n\},$$
$$A_n = a[\alpha_n].$$

\E@ Tal2.6.1 <2>

Then we have $|E_n| = 2^{n-1}$ and $|J_n| = 2^n - 1$. The reason for this particular choice of the α_n 's will only become evident when bounds involving those constants are studied. The following calculation is an example.

First note that $||a||^2 = \sum_{n \in \mathbb{N}} \sum_{i \in E_n} a_i^2$ and, by monotonicity of the $(a_i : i \in \mathbb{N})$ sequence we have $A_n \ge a_i \ge A_{n+1}$ for $i \in E_n$. Hence

UPPER :=
$$\sum_{n \in \mathbb{N}} |E_n| A_n^2 \ge ||a||^2 \ge \sum_{n \in \mathbb{N}} |E_n| A_{n+1}^2$$

= $\frac{1}{2} (\text{UPPER} - |E_1| A_1^2).$

From the fact that $|E_1|A_1^2 = a_1^2 \le ||a||^2$ it then follows that



Remark. This form of "condensation argument" is often attributed (Rudin, 1976, Theorem 3.27) to Cauchy.

The argument leading to $\langle 2 \rangle$ works separately for each t in T, using a sequence (as described in Lemma $\langle 6 \rangle$) of non-negative integers $p_0 = 0, p_1, p_2, \ldots$ depending on t. The construction works a block at a time, starting from

$$R_0 := t = (\tau_1, \tau_2, \dots) \qquad \text{where } \tau_k := (t_i : i \in E_k) \in \mathbb{R}^{E_k}.$$

It helps to rescale the problem so that $||a||_2 = 1$. Then the fact that $\sum_i t_i^2/a_i^2 \leq 1$ shows that not only do we have $t_i^2 \leq a_i^2$ for each *i* but also



$$1 = \sum_{i \in \mathbb{N}} a_i^2 \ge \sum_{i \in \mathbb{N}} t_i^2 = \sum_{n \in \mathbb{N}} |\tau_n|_2^2 \quad \text{and} \\ 1 \ge \sum_{i \in \mathbb{N}} t_i^2 / a_i^2 \ge \sum_{n \in \mathbb{N}} |\tau_n|_2^2 / A_n^2.$$

The argument leading to $\langle 2 \rangle$ is recursive. The first step replaces $R_0[J_1] := (R_0[i] : i \in J_1)$ using an x_1 in \mathbb{R}^{J_1} such that the difference $r_1 := R_0[J_1] - x_1$ has $|r_1|_2 \leq 2^{-p_1/2}$. The working vector is then

$$R_1 := (r_1, \tau_2, \tau_3, \dots).$$

The next step finds an x_2 in \mathbb{R}^{J_2} such that the difference $r_2 := R_1[J_2] - x_2$ has $|r_2|_2 \leq 2^{-p_2/2}$, with new working vector

$$R_2 := (r_2, \tau_3, \tau_4, \dots)$$

And so on.

Each x_k in \mathbb{R}^{J_k} defines an element $\widetilde{x}_k := (x_k, 0, 0, \dots)$ of ℓ^2 . For each t in T the challenge is to choose x_k from a subset V_k of \mathbb{R}^{J_k} with $|V_k|$ suitably small size. The approximating set U_n then consists of all possible sums $\widetilde{x}_1 + \cdots + \widetilde{x}_n$ with $\widetilde{x}_k \in V_k$ as t ranges over T.

Remark. The tildes are just to remind you that we need to identify the x_k from \mathbb{R}^{J_k} with an element of ℓ^2 . I was rather tempted to abuse notation by writing x_k instead of \tilde{x}_k .

\E0 norm.bnds $<\!3\!>$

The next Lemma establishes the important properties of the p_n , which I initially write as p[n, t]'s to emphasize the dependence on t.

Lemma. For a fixed $t = (\tau_1, \tau_2, ...)$ in $T := \mathcal{E}[a]$ with $||a||_2 = 1$ define

$$\beta[n,t] := \sup\{k \in \mathbb{N}_0 : 2^{-k} \ge |\tau_n|_2^2\}$$

and then recursively define a sequence of nonnegative integers p[n,t] by $p[0,t] := p_0 := 0$ and

$$p[n,t] := \min(p[n-1,t]+2,\beta[n,t])$$
 for $n \ge 1$.

Then:

 $<\!6\!>$

pn

- (i) $p[n,t] \leq p[n-1,t] + 2$ for $n \geq 1$ so that $p[n,t] \leq 2n$ for each n and $t \in T$.
- (ii) $2^{-p[n,t]} \ge 2^{-\beta[n,t]} \ge |\tau_n|_2^2$ for each n.
- (iii) $\sum_{n \in \mathbb{N}_0} 2^{n/2 p[n,t]/2} \leq c_1$ where c_1 is a universal constant.

Proof. As the argument works independently for each t let me drop the tfrom the notation, abbreviating p[n, t] to p_n and so on. Also I'll write h_n for $|\tau_n|^2$.

Assertions (i) and (ii) come directly from the definition of p_n as a minimum of two quantities. Assertion (iii) is a bit more subtle.

Define $\mathbb{N}_{\beta} := \{n \in \mathbb{N} : p_n = \beta_n\}$. If $n \in \mathbb{N}_{\beta}$ then β_n is finite and

$$2^{-p_n} = 2^{-\beta_n} \ge h_n > 2^{-\beta_n - 1} = 2^{-p_n - 1}$$
 implying $2h_n > 2^{-p_n}$.

Consequently,

$$\sum_{n \in \mathbb{N}_{\beta}} 2^{-p_n} / A_n^2 \le \sum_{n \in \mathbb{N}_{\beta}} 2h_n / A_n^2 \le 2 \qquad \text{by inequality } <5>.$$

By CAUCHY-SCHWARZ and inequalities $\langle 3 \rangle$ and $\langle 8 \rangle$ we then have

$$\left(\sum_{n\in\mathbb{N}_{\beta}} 2^{n/2-p_n/2}\right)^2 \le \sum_{n\in\mathbb{N}_{\beta}} \frac{2^{n-p_n}}{2^n A_n^2} \sum_{n\in\mathbb{N}_{\beta}} 2^n A_n^2 \le 6.$$

If $\mathbb{N}\setminus\mathbb{N}_{\beta}$ is nonempty then it must be union of disjoint stretches I_1, I_2, \ldots of consecutive integers. Suppose

$$I := \{k + j : 1 \le j \le \ell\} \subset \mathbb{N}_0 \setminus \mathbb{N}_\beta$$

for some $k \in \{0\} \cup \mathbb{N}_{\beta}$ and some $\ell \in \mathbb{N}$. By assumption $p_n < \beta_n$, which forces $p_n = 2 + p_{n-1}$, for $n \in I$. It then follows that $p_{k+j} = p_k + 2j$ for $1 \leq j \leq \ell$ so that

$$\sum_{n \in I} 2^{n/2 - p_n/2} = 2^{k/2 - p_k/2} \sum_{j=1}^{\ell} 2^{-j/2} \le c_0 2^{k/2 - p_k/2},$$

where $c_0 = (\sqrt{2} - 1)^{-1} \approx 2.5$. A similar inequality holds if |I| is infinite. If k = 0 the upper bound equals c_0 . If $k \in \mathbb{N}_\beta$ then $2^{k/2-p_k/2}$ is one of the summands on the left-hand side of inequality $\langle 9 \rangle$. Assertion (iii) follows.

< 7 >\E@ beta.finite

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< 9 >
\E@ n2p2
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\EC sum.NNbeta $<\!\!8\!\!>$

3 The main construction: approximation for ellipsoids

S:approximation

Throughout this Section I assume $||a||_2 = 1$ so that by inequalities $\langle 4 \rangle$ and $\langle 5 \rangle$,

$$\max\left(\sum_{n\in\mathbb{N}}|\tau_n|_2^2,\sum_{n\in\mathbb{N}}|\tau_n|_2^2/A_n^2\right)\leq 1\qquad\text{for each }t=(\tau_1,\tau_2,\dots)\text{ in }T.$$

Define $\delta(p) := 2^{-p/2}$ for $p = 0, 1, \dots$ From Lemma 6 with $p_n = p[n, t]$ we have $|\tau_n|_2 \leq \delta(p_n)$ and $p_n \leq p_{n-1} + 2$ for $n \geq 1$, implying

$$\delta(p_n) = 2^{-p_n/2} \ge 2^{-p_{n-1}-1} = \frac{1}{2}\delta(p_{n-1})$$

For each finite subset J of \mathbb{N} and each non-negative integer p define

$$B(J,p) := \{ y \in \mathbb{R}^J : |y|_2 \le 4\delta(p) \}$$

the euclidean ball with center 0 and radius $4 \times 2^{-p/2}$. Let V(J,p) be a $\delta(p)$ packing set for B(J,p). From PTTM SECTION 10.4

$$4^{|J|} \le |V(J,p)| := \operatorname{Pack}(\delta(p), B(J,p)) \le 12^{|J|}$$

Each point of B(J,p) lies within distance $2^{-p/2}$ of V(J,p). Thus there is a map $\Psi_{J,p}: B(n,p) \to V(J,p)$ for which

$$|y - \Psi_{J,p}y|_2 \le 2^{-p}$$
 for each y in $B(J,p)$.

Remark. Without loss of generality it can be assumed that $0 \in V(J, p)$ for all (J, p), which ensures that $\Psi_{J,p} 0 = 0$.

1: procedure APPROX(t)% fixed $t = (\tau_1, \tau_2, ...)$ in T Initialize: $n \leftarrow 1$; $R_0 \leftarrow t$. 2: while $R_{n-1} \neq 0$ do 3: 4: $y_n \leftarrow R_{n-1}[J_n].$ $x_n \leftarrow \Psi_{J_n,p_n}(y_n).$ 5: $R_n \leftarrow (r_n, \tau_{n+1}, \tau_{n+2}, \dots)$ where $r_n := y_n - x_n$. 6: Increment n by 1. 7: end while 8: 9: end procedure It helped me to run through a few steps of the while loop.

For
$$n = 1$$
:

$$|y_1|_2 = |\tau_1|_2 \le \delta(p_1) \text{ so that } y_1 \in B(J_1, p_1),$$

$$x_1 = \Psi_{J_1, p_1} y_1 \in V(J_1, p_1) \quad \text{AND} \quad |r_1|_2 \le \delta(p_1)$$

$$R_1 = (r_1, \tau_2, \tau_3, \dots).$$

\E@ ball.pack <11>

\E@ delpn <10>

For n = 2:

$$|y_2|_2 = |(r_1, \tau_2)|_2 \le |r_1|_2 + |\tau_2|_2 \le \delta(p_1) + \delta(p_2) \le 4\delta(p_2)$$

so that $y_2 \in B(J_2, p_2)$,
 $x_2 = \Psi_{J_2, p_2}^{-1} y_2 \in V_{J_2, p_2}$ and $|r_2|_2 \le \delta(p_2)$,
 $R_2 = (r_2, \tau_3, \tau_4, \dots)$.

For n = 3:

$$\begin{aligned} |y_3|_2 &= |(r_2, \tau_3)|_2 \le |r_2|_2 + |\tau_3|_2 \le \delta(p_2) + \delta(p_3) \le 4\delta(p_3) \\ \text{so that } y_3 \in B(J_3, p_3), \\ x_3 &= \Psi_{J_3, p_3}^{-1} y_3 \in V_{J_3, p_3} \text{ and } |r_3|_2 \le \delta(p_3), \\ R_3 &= (r_3, \tau_3, \tau_5, \dots). \end{aligned}$$

In general we have $y_n \in B(J_n, p_n)$ and $x_n \in V(J_n, p_n)$ with that $|r_n|_2 \le 4\delta(p_n)$ with

$$R_n = (r_n, \tau_{n+1}, \tau_{n+2}, \dots) = (y_n - x_n, \tau_{n+1}, \tau_{n+2}, \dots)$$
$$R_{n-1} = (r_{n-1}, \tau_n, \tau_{n+1}, \dots) = (y_n, \tau_{n+1}, \dots).$$

so that $R_{n-1} - R_n = (x_n, 0, 0, \dots) =: \widetilde{x}_n$. It follows that

$$t - R_n = \sum_{k=1}^n \left(R_{n-1} - R_n \right) = \sum_{k=1}^n \widetilde{x}_n =: u_n$$

The remainders tend to zero as n goes to infinity because

$$||R_n||_2 \le |r_n|_2 + \sum_{k>n} |\tau_k|_2 \le 4 \sum_{k\ge n} \delta(p_k).$$

Finally

$$\sum_{n \in \mathbb{N}} 2^{n/2} d(t, U_n) \leq \sum_{n \in \mathbb{N}} 2^{n/2} ||t - u_n||_2 = \sum_{n \in \mathbb{N}} 2^{n/2} ||R_n||_2$$

$$\leq 3 \sum_{n \in \mathbb{N}} 2^{n/2} \sum_{k \in \mathbb{N}} \{k \leq n\} 2^{-p_k/2}$$

$$= 3 \sum_{k \in \mathbb{N}} 2^{-p_k/2} \sum_{n \in \mathbb{N}} 2^{n/2} \{k \leq n\}$$

$$\leq 3 \sum_{k \in \mathbb{N}} 2^{-p_k/2} 2^{(k+1)/2} / (\sqrt{2} - 1),$$

which by Lemma $\langle 6 \rangle$ is smaller than $3c_1\sqrt{2}/(\sqrt{2}-1)$.

As t ranges over T we get different $\{p[n,t]\}$ sequences, but always in the range $1 \le p_k \le 2k$ and

$$y_k \in V_k := \bigcup_{1 \le p \le 2k} V(J_k, p).$$

By inequality <11>,

$$|V_k| \le \sum_{1 \le p \le 2k} 12^{|J_k|} \le 2k \times 12^{2^k} \le 12^{2^{k+1}}$$

Similarly, $|U_n| \leq \prod_{k=1}^n |V_k|$, implying

$$\log_2 |U_n| \le \sum_{k=1}^n 2^{k+C_1} \le 2^{k+C_2},$$

for some universal constants C_1 and C_2 .

S:packing

4

Why chaining with packing numbers can fail

If $a \in \ell^2$ then $\langle 3 \rangle$ gives $\sum_{n \in \mathbb{N}} 2^n A_n^2 < \infty$, where $A_n := a[2^{n-1}]$. However it does not imply finiteness of $\sum_{n \in \mathbb{N}} 2^{n/2} (A_n - A_{n+1})$, as shown by the example $a_i := (\sqrt{i} \log_2(2i))^{-1}$ where $A_n/2 = 2^{-n/2} n^{-1}$.

For such an *a* the packing bound gives $\mathbb{P}_{t \in \mathcal{E}[a]} X_t \leq \infty$.

Suppose F is an ϵ -packing set for $T := \hat{\mathcal{E}}[a]$ with |F| = N. Define $B_k := \{s \in \mathbb{R}^{E_k} : |s|_2 \leq A_k\}$. Consider the vector $t = \tilde{s} = (0, \ldots, 0, s, 0, \ldots)$. By definition of a packing set, there exists an f in F for which

$$|s - f[E_k]|_2 \le ||t - f||_2 \le \epsilon$$

Consequently, the set of vectors $\{f[E_k] : f \in F\}$ is an ϵ -covering set for B_k . Inequality $\langle 11 \rangle$ gives $N \geq (A_k/\epsilon)^{|E_k|}$.

$$N \ge \left(A_k/\epsilon\right)^{|E_k|}.$$

In particular, it follows that

$$\log_2 \operatorname{Pack}(\epsilon, T) \ge |E_k| = 2^{k-1} \quad \text{if } \epsilon \le A_k/2$$

and hence

$$\int_{A_{k+1/2}}^{A_k/2} \sqrt{\log \text{PACK}(r,T)} \, dr \ge \text{const.} 2^{k/2} \, (A_k - A_{k+1}) \ge \text{const.}/k.$$

Sum over k to deduce that

$$\int_0^{\operatorname{diam}(T)} \sqrt{\log \operatorname{pack}(r,T)} \, dr = +\infty$$

Remark. Check §10.4 to see if pack vs. cover makes a difference.

Problems

S:Problems

 $\mathbf{5}$

[1]

P:special.t

Suppose $t \in \mathcal{E}[a]$ with $t_i = 0$ for $i \notin E_k$. Assume $0 \in V(J, p)$ for all (J, p), which ensures that $\Psi_{J,p}0 = 0$. Find the sequences p[n, t] and the vectors R_n generated by procedure APPROX. (Note the dependence on k.)

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